### An Introduction to Numeric Planning Representation and Search Algorithms

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### Goals and Assumptions of this Lecture

Introduction to the field of Numeric AI Planning in which we

- Define the problem, present the main representation language, main planning algorithms and systems
- Focus on the state-based model for (numeric) planning with
  - discrete time and "instantaneous" actions
  - full observability of states
  - deterministic action effects
  - no state changes other than those caused by the plan

State-based Numeric Planning

Representation and Languages (PDDL2.1)

Solving Numeric Planning through Heuristic Search

Relaxation-based Heuristics for Numeric Planning

### Basic State-based Model (Classical Planning)

Plan generation as state-transition problem:

- a finite and discrete state space S (potentially huge!)
- a known <u>initial state</u>  $s_0 \in S$
- a set  $S_G \subseteq S$  of goal states
- a finite set of <u>actions</u> A
- actions  $A(s) \subseteq A$  applicable in each  $s \in S$
- a deterministic state transition function  $s' = \gamma(a, s)$  for  $a \in A(s)$
- action <u>costs</u>  $c(a, s) \ge 0$

A solution is a sequence of applicable actions that maps  $s_0$  into  $S_G$ It is optimal if minimizes the sum of action costs

Related problems: plan existence, plan validation, plan revision, goal recognition, interleaved planning/execution

### Propositional Planning Language

In a propositional planning language

- states represented by propositions (atoms, Boolean vars, facts)
- $\gamma$  and A represented by action preconditions/effects over atoms

A problem in the STRIPS language is a tuple  $\langle F, A, I, G \rangle$ :

- F = set of all atoms
- A = set of all actions
- $I \subseteq F$  is an initial state Note: all other states are *implicit*!
- $G \subseteq F$  is a set of goals

Each  $a \in A$  represented by

- Positive effects Add(a) ⊆ F and negative effects Del(a) ⊆ F
- Preconditions Pre(a) ⊆ F

Many extensions: conditional effects, non-atomic preconditions, domain axioms, state trajectory constraints, etc.

Numbers are very useful for modelling problems involving, e.g.,

- Available Resources (fuel, energy, money, ...)
- *Physical quantities* (temperature, velocity, pressure, density, ...)
- Spatial information (cartesian coordinates, GPS position, size, ...)
- $\Rightarrow$  We add numbers in the planning model and in the language:
  - Actions with numeric preconditions and numeric effects fuel-level > 100 before moving and reduced by 88.5 after moving

#### • Numeric goals

target location has been visited and fuel-level > 10

States are specified by variables whose values are (rational) numbers

- an infinite space of numeric states S
- a known initial **numeric** state  $s_0 \in S$
- a **possibly infinite** set  $S_G \subseteq S$  of goal states
- A finite set of actions A
- actions  $A(s) \subseteq A$  applicable in each  $s \in S$
- a deterministic state transition function  $s' = \gamma(a, s)$  for  $a \in A(s)$
- action costs  $c(a, s) \ge 0$

#### Note: $\gamma$ has a potentially **infinite number of state transitions**

 $\Rightarrow$  How do we represent states,  $\gamma$ , and action preconditions/effects?

**PDDL** (Planning Domain Definition Language) is the mostly used standard language for planning

Incrementally developed

- PDDL1 [McDermott and others, 1998]
- PDDL1.2 [Bacchus, 2000]
- PDDL2.1 [Fox and Long, 2003]
- PDDL2.2 [Hoffmann and Edelkamp, 2004]
- PDDL+ [Fox and Long, 2006]
- PDDL3.0 [Gerevini and Long, 2006]
- PDDL3.1 [Helmert et al., 2008]

PDDL2.1 introduced numeric planning and temporal planning in PDDL

PDDL+ and PDDL3: numeric planning problems with *events*, *processes*, *soft goals*, *metric state-trajectory constraints* 

### Numeric Planning Supported in PDDL2.1

A planning problem is a tuple  $\langle F, X, A, s_0, G \rangle$  where

- F is a finite set of **Boolean variables** (propositions/atoms/facts)
- X is a set numeric variables over  $\mathbb{Q}$  (numeric fluents)
- A is a set of **actions** (with preconditions and effects)
- $s_0$  is the **initial problem state** (an assignment to all vars in  $F \cup X$ )
- *G* are the **problem goals** (subset of atoms in *F* plus numeric conditions over the vars in *X*)
- ⇒ Which class of action models (preconditions and effects) and goals can be represented in PDDL2.1?

### Action Precondions/Effects & Goals (PDDL2.1 Models)

#### **Preconditions** (*Pre*(*a*)):

- Truth conditions on Boolean vars
- Conditions on numeric vars of the form "ξ Rel k"
   k ∈ Q, ξ is arithmetic expression, Rel ∈ {≤, <, =, >, ≥}

#### **Effects** (*Eff*(*a*)):

- Truth changes to Boolean vars
- Changes to numeric vars of the form *increase, decrease, assign, scale-up/down* by an *arithmetic expression*
- $\Rightarrow$  Should not conflict! (e.g. *decrease* and *scale up* the same var)

#### Problem goals:

- Propositional goals as in propositional planning
- Numeric goals (same as the numeric action conditions)

Note that

In classical planning (STRIPS) action **effects are state-independent**: every action has predefined effects that are the same for every state where the action can be applied

But this does not hold for numeric effects

E.g.: (decrease fuel 10) assigns 40 to fuel if it was 50, 0 if it was 10

 $\Rightarrow$  The "state dependency" of the action effects can make the design of effective algorithms for numeric planning more difficult

(see heuristic search techniques in the 2nd part of this lecture)

### Example in PDDL Syntax (plant watering)

**Task**: one or more agents on a grid-like map have to water some plans at different locations using water loaded and carried from sources (taps) available at different locations

- water tank of an agent has limited capacity
- water loaded in the tank by smaller container (multiple load actions)
- each plant needs its amount of water
- each move (up, down, left, right) and water load/unload has a cost



## Example in PDDL (objects, predicates and functions)

PDDL supports the definition of

• Problem objects with types

(:objects agent1 agent2 ... - agent)

<u>Predicates</u>

(:predicates (open ?t - wathersource))

Grounded predicate (open tap1) is a Boolean var/proposition

• *Functions* (numerical fluents/numeric vars)

(:functions (x-pos ?a - agent)) (y-pos ?a - agent))

(x-pos agent1) and (y-pos agent2) are numeric fluents (vars)

• Action models with parameters (next slide)

 $\label{eq:problem action} \textbf{Problem action} = \text{PDDL} \text{ action with instantiated parameters}$ 

If ?a = agent1 and ?t = watercourse, the involved numeric vars are:
 (x-pos agent1) (y-pos agent1) (carrying) (max\_load) (total\_loaded)

#### Applicable action

- An action *a* is applicable in a state *s* when  $s \models Pre(a)$
- An action a applied in s generates the state s' where the variable values of s are changed according to Eff(a)

#### Applicable action sequence

A sequence of actions  $\langle a_0, ..., a_{n-1} \rangle$  is a **applicable in state**  $s_0$  when

- a<sub>0</sub> is applicable in s<sub>0</sub>
- each action a<sub>i</sub> is applicable in the state s<sub>i</sub> generated by the application of a<sub>i-1</sub> in s<sub>i-1</sub>

**Valid plan (solution)** is a sequence of actions  $\langle a_0, ..., a_{n-1} \rangle$  that, when applied in  $s_0$ , generates a sequence of states  $\langle s_1, ..., s_n \rangle$  such that  $s_n \models G$ .

### Non-Linear Plans in PDDL2.1 (informally)

A valid plan may also be a sequence of <u>sets</u> of actions  $\langle A_0, ..., A_{m-1} \rangle$  where each pair of actions in  $A_i$  does not interfere (i = 0...m - 1)

**Interference of** *a* with *b* (informally)

• An effect in Eff(a) modifies the value of a var involved in Pre(b)

- (decrease fuel 10)  $\in$  Eff(a) and (fuel > 100)  $\in$  Pre(b)

• The effects in Eff(a) do not commute with the effects in Eff(b)

- (assign x 0)  $\in$  Eff(a) and (assign y x)  $\in$  Eff(b)

### Plan Metrics for Numeric Planning in PDDL2.1

PDDL plans have a **plan quality** defined through numeric variables

- Action costs can be defined in terms of a numeric effect increasing var total-cost
- Quality of plans can be defined in terms of
  - minimizing total-cost or
  - minimizing a numeric (arithmetic) expression
- Example using simplified PDDL syntax:

(:metric minimize (+ moving\_cost (\* 2 total\_loaded)))

Actions can only increase moving\_cost and total\_loaded

Numeric planning with PDDL2.1 is **semi-decidable** [Helmert 2002] (harder than classical planning)

- In numeric planning plans can have no bound in their length
- In principle we can find a plan by enumerating all sequences of actions with increasing length
- But enumeration may not terminate if a plan does not exists (semi-decidability)

**Decidable** for some other fragments of PDDL2.1 <sup>(2)</sup> [see Helmert 2002]

Researchers have studied numeric planning with one or more of the following restrictions to PDDL2.1:

- The numeric expression  $\xi$  in a precondition  $\xi \operatorname{Rel} k$  is **linear**
- assign effects excluded
- *scale-up/down* effects **excluded**
- *increase/decrease* effects modify the vars by **constants** instead of generic arithmetic expressions- E.g. (increase x 5)

A Simple Numeric Problem has all these but still semi-decidable  $\bigcirc$ 

(see how to solve them in the 2nd part of the lecture)

### Some Limits and Extensions of PDDL2.1

- A numeric var may be undefined in a state (no initial value)
   With x is undefined and y = 5, are these true/false/undefined?
   (x > 0) (not (x > 0)) (or (x > 0) (y < 10))</li>
- $\bigcirc$  Only arithmetic expressions: no *min*, *max*, *log*, *exp*, *sin*, *cos*, etc.
- O Conflicts of numeric effects of type scale-up/down could be removed
- ③ Numeric parameters for an action are forbidden (finite action set)
- Nevertheless PDDL2.1 is a powerful language and researchers have addressed some of these limits [e.g., Savas-et-al 2016, Scala-et-al, 2016]
- Wumeric planning can be combined with temporal planning (PDDL2.1/2.2) and enriched with: state trajectory constraints (PDDL3), exogenous events and continuous processes (PDDL+)

 $\Rightarrow$  The full PDDL language is very powerful!



Given (1) a set of types of goods (2) a set of markets ( $\mathbf{M}$ ) selling different types and amounts of goods at different prices, (3) a demand of each type of goods to be purchased and transported to some depot ( $\mathbf{D}$ ),

 $\Rightarrow$  satisfy the demand minimizing the routing cost of the trucks and the purchasing cost

6 different PDDL versions (see IPC5) with simplifications and extensions

```
(:predicates (at ?t - truck ?p - place))
```

```
(:functions (on-sale ?g - goods ?m - market)
      (drive-cost ?p1 ?p2 - place)
      (price ?g - goods ?m - market)
      (bought ?g - goods)
      (request ?g - goods)
      (total-cost))
```

(:objects

market1 market2 ... - market depot1 depot2 ... - depot truck1 track2 ... - truck goods1 goods2 ... - goods)

```
(:init (= (price goods1 market1) 15)
       (= (price goods1 market2) 8)
       . . . .
       (= (on-sale goods1 market1) 4)
       (= (on-sale goods1 market2) 9)
       . . . .
       (at truck1 depot0)
       (at truck2 depot1)
       . . . .
       (= (drive-cost depot0 market1) 381.20)
       (= (drive-cost market1 market2) 1033.70)
       . . . .
       (= (bought goods1) 0)
       (= (bought goods1) 0)
       . . . .
       (= (request goods1) 20)
       (= (request goods2) 23)
       . . . .
       (= (total-cost) 0))
(:goal (and (>= (bought goods1) (request goods1) ....)))
(:metric minimize (total-cost))
```



### More on Numeric Planning in PDDL?

- Many benchmarks from the International Planning Competitions https://www.icaps-conference.org/competitions/
- PDDL Wikipedia page: https://en.wikipedia.org/wiki/Planning\_Domain\_Definition\_Language
- Paper on PDDL 2.1 [Fox and Long, JAIR 2003]: https://www.jair.org/index.php/jair/article/view/10352
- Paper on PDDL3.0 [Gerevini et al., AIJ 2009]: https://www.sciencedirect.com/science/article/pii/S0004370208001847
- Complexity of PDDL Numeric Planning [Helmert, AIPS 2002]
- Book on PDDL: "An Introduction to the Planning Domain Definition Language", Haslum et al., 2019, Morgan & Claypool Publishers

### Solving Numeric Planning Problems: Approaches

As for propositional planning, several different approaches and algorithms

• **Planning by compilation**: Translates the original problem into an equivalent one solved by existing solvers

E.g., SAT for propositional planning, SMT for numeric planning

• **Planning as heuristic search**: We use general search algorithms with effective *domain-independent heuristics* 

E.g., Best-first search, A\*, local search

- directional (forward/backward) search in the space of states E.g., the Metric-FF planner
- non-directional search in the space of partial plans E.g., the LPG planner

In this lecture we focus on forward state-based search

#### Informal Definition

State space search is about iterating over paths of a transition system until either

- there is no new path to explore
- a path proved to be a valid solution is found

Forward State Space Search is the variant where we always start from some initial condition

Forward State Space Search for Planning revolves around a few, simple ideas:

- Map problem into underlying transition system (TS) and seek trajectory belonging to it. Two states  $s_0$  and  $s_1$  are connected whenever there is an action applicable in  $s_0$  yielding  $s_1$
- Focus on reachable sequences only! Simulate possible futures starting from initial condition
- *Discard Loops!* Focus only on simple plans; each state does not need to be visited more than once
- Do not construct the whole TS upfront, but incrementally. Realistic planning instances induce way too large TSs

Best-First-Search(problem)

 $\label{eq:Q} \begin{array}{l} \mathsf{Q} = \{ \text{ Node(problem.init, 0, h(problem.init))} \} \\ \textbf{while } \mathsf{Q} \text{ is not empty } \textbf{do} \\ \texttt{n} = \textbf{pop\_best\_node}(\mathbf{Q}) \\ \textbf{if } \texttt{n.s is a goal state then} \\ \textbf{return extract\_plan(n.s)} \\ \textbf{for all } \texttt{s'} \in \textbf{succ(n.s) do} \\ \texttt{insert}(\mathsf{Q}, \texttt{Node}(\texttt{s', n.g+1, h(s')})) \\ \textbf{return problem is not solvable} \end{array}$ 

- Q contains all plan-prefixes, each implicitly represented by a node with a pointer to its parent
- BFS operates Q in two ways:
  - pop next node according to some ordering function f
  - push new plan prefixes by appending new nodes obtained through successor function

For simplicity, duplicate checking and node re-opening is not detailed

#### Best-First-Search(problem)

 $\begin{array}{l} Q = \{ \mbox{ Node(problem.init, 0, h(problem.init))} \} \\ \mbox{while } Q \mbox{ is not empty } \mbox{do} \\ n = \mbox{pop_best_node}(Q) \\ \mbox{if } n.s \mbox{ is a goal state then} \\ \mbox{ return extract_plan}(n.s) \\ \mbox{ for all } s' \in \mbox{succ}(n.s) \mbox{ do} \\ \mbox{ insert}(Q, \mbox{ Node}(s', n.g+1, h(s'))) \\ \mbox{return problem is not solvable} \end{array}$ 

• f determines the particular algorithm

• UCS 
$$-> f(n) = g(n)$$

• 
$$A^* \to f(n) = g(n) + h(n)$$

• 
$$WA^* \to f(n) = g(n) + w \times h(n)$$

• GBFS -> f(n) = h(n)

#### For simplicity, duplicate checking and node re-opening is not detailed

Solving numeric planning through forward state space search is straightforward:

- State representation. Variables that keep track of values of numeric fluents
- 2 Extend successor function using numeric planning semantics
  - An action is applicable in state s iff its precondition evaluates to true in s. This needs to include reasoning over numeric conditions (e.g., yx + z > 5 is satisfied in s)
  - Successor of a state computed using actions' effects semantics
- 3 Goal test. Ability to evaluate formulas involving numeric terms

#### Best-First-Search(problem)

 $\label{eq:Q} \begin{aligned} & Q = \{ \text{ Node(problem.init, 0, h(problem.init))} \} \\ & \textbf{while } Q \text{ is not empty do} \\ & n = \textbf{pop\_best\_node(Q)} \\ & \textbf{if } n.s \text{ is a goal state then} \\ & \textbf{return extract\_plan(n.s)} \\ & \textbf{for all } s' \in \textbf{succ(n.s) do} \\ & \text{insert}(Q, \text{ Node(s', n.g+1, h(s'))}) \\ & \textbf{return problem is not solvable} \end{aligned}$ 

**goal state check**, **succ** are the key points that need to be revised in order for getting the schema work for numeric planning problems

For simplicity, duplicate checking and node re-opening is not detailed

### Plant-Watering Example Search Tree



- Actions: East, West, South, North, Load-Water, Pour-Plant
- Goal: Get plant poured



- Number of iterations (AKA size of the search tree) can be prohibitively large. Note that numeric planning transition system is infinite.
- The function *f* plays critical role. It is what guides the exploration
- Even though we cannot change the action costs, we can change the way we predict how distant the goal is and use an informed BFS (e.g., A\*, GBFS, WA\*)
- That is: We can change the heuristic

Known facts:

- If heuristic is perfect (true distance to the goal for each state), number of expansions is **linear with the length of the solution**
- The closer the heuristic is to the perfect estimate (called *h*<sup>\*</sup>), the less is the number of expanded nodes (Dechter and Pearl 1983)
- Most research look into making heuristic functions as much precise as possible

Almost perfect heuristics can still cause exponential blow ups, Helmert and Roger 2008.

### Definition (Heuristic)

Heuristic is a function  $h: S \rightarrow Q \cup \{\infty\}$ 

### Definition (Admissibility)

Heuristic *h* is said to be admissible if  $\forall s \in S.h(s) \leq h^*(s)$ ; A\* with admissible heuristic finds optimal solutions.

#### Definition (Safe-Pruning)

Heuristic *h* is said to be safe-pruning if  $\forall s \in S.h(s) = \infty \Rightarrow$  there is no solution from *s* 

#### Definition (Tractability)

Heuristic is tractable if computable in poly-time (w.r.t. problem size). Heuristic should be easy than the problem they are approximating

### How to Define Heuristics for Numeric Planning Problems

- The user defines it. Can work well, but how general can it ever be?
- Construct an algorithm that can look at the problem and tell us what the value of the heuristic is. Automatically!
- Heuristics of this kind are often called Domain Independent Heuristics

We will show domain independent heuristics for numeric planning, by investigating their properties. Admissibility, Safe-Pruning, Tractability

### Relaxation-Based Heuristics. The Intuition



- Devise a new problem that relaxes the problem you are interested to solve
- Make sure the arising relaxed problem is easier, yet contains enough information
- Solve the relaxed problem, and use the solution to guide the solution of the new problem. This is what gives you your heuristic

Bonet and Geffner 2001, Hoffmann and Nebel 2001

### Relaxation-based Approach in Numeric Planning

- Interval-based Relaxation (Hoffmann 2003, Aldinger et al. 2015)
  - Decompose/ignore interactions among the numeric variables of your problem
  - Characterized by starting from initial state up to a point in which the goal is relaxed satisfied
- Subgoaling-based Relaxation (Bonet and Geffner 2001, Scala et al. 2016)
  - Decompose/ignore interactions among the numeric subgoals of your problem
  - Characterized by starting regressively from the goal condition up to the point in which all action supporters have been found

Both frameworks can be used to devise a variety of heuristics

### Interval-based Relaxation (Intuitively)

• Map numeric state variables into intervals (many states at once)

• Ex. 
$$s: \{x = 0, y = 0\} \rightarrow s^+: \{x = [0, 0], y = [0, 0]\}$$

- Change semantics of effects through convex union (intervals only grow)
  - Ex.  $s^+[x+=1] = \{x = [0,1], y = [0,0]\}$  Just update the max
  - Ex.  $s^+[x = 5] = \{x = [0, 5], y = [0, 0]\}$  Intervals don't have holes!
- Change semantics of formulas so as to over-approximate satisfiability
  - Ex. x = 10 is satisfied in [0, 10]
  - Ex. x + y = 10 ?

Each operation is computed so as to enclose all values that can ever be attained

#### Arithmetical Operations

• 
$$x + y = [\underline{x} + \underline{y}, \overline{x} + \overline{y}];$$

- $x y = [\underline{x} \overline{y}, \overline{x} \underline{y}];$
- $x \times y = [min(\underline{xy}, \underline{xy}, \overline{xy}, \overline{xy}), max(\underline{xy}, \underline{xy}, \overline{xy})];$
- $x \div y = [min(\underline{x} \div \underline{y}, \underline{x} \div \overline{y}, \overline{x} \div \underline{y}, \overline{x} \div \overline{y}), max(\underline{x} \div \underline{y}, \underline{x} \div \overline{y}, \overline{x} \div \underline{y}, \overline{x} \div \overline{y})]$ (if  $0 \notin y$  otherwise one of the bounds diverges).

#### Note

Interval analysis is a very powerful framework, Many other mathematical operations can be defined. More dtails in Moore et al. 2009, and their application into planning in Aldinger et al. 2015, Scala et al. 2016.

#### Example 1

Condition  $\psi : x + y = 0$ ; Interval State  $s : \{x = [0, 10], y = [-15, -15]\}$ Is condition satisfied? No, indeed [0, 10] + [-15, -15] = [-15, -5]

#### Example 2

Condition  $\psi : x + y = 0$ ; Interval State  $s : \{x = [0, 10], y = [-15, -5]\}$  Is condition satisfied? Yes, indeed [0, 10] + [-15, -5] = [-15, 5] and it is possible to pick 0 within [-15, 5] so as to satisfy the condition

### Interval-based Relaxation. The problem

#### Interval-based Relaxation

Given a planning problem  $\Pi = \langle s_0, A, G, X \rangle$ ,  $\Pi^+ = \langle s_0^+, A^+, G^+, X^+ \rangle$  is its Interval-based relaxation where

- $s_0^+$  is represented as a set of intervals.
- A<sup>+</sup> syntactical equivalent to A, but different semantics:
  - 1 Successor state via convex union
  - **2** Preconditions only relaxed satisfied  $(\exists v \in [\underline{x}, \overline{x}] : v \models pre(a))$
- $G^+$  is syntactical equivalent to G. Semantics as for action precondition
- $X^+ \equiv X$ . Variables don't change, but expressions follow Interval Algebra

#### How to solve it?

Two step process:

- **1** Reachability analysis of the goal (forward analysis)
- 2 Extraction of a relaxed plan solution (backward analysis)

### Reachability through Example



#### The process

- Determine action applicability using intervals
- Apply actions with convex union up to the point where goal is satisfied
- Store actions for each encountered level

#### Observations

- **1** Conditions satisfied at level *i* will also be satisfied in any level j > i
- 2 Applicable actions never decrease

# All good when the problem is solvable **What about when the problem** isn't solvable?

- Metric-ff **mneed criteria** (Hoffman 2003) (linear expressions, acyclic assignments)
  - Compute max sufficient value for each variable. Stop when fix point is reached, or all intervals contain mneed
- Asymptotic reachability (Aldinger et al. 2015, Scala et al. 2016) (linear, non-linear, state dependent effects that can be acyclic, too)
  - make hidden state dependency explicit in preconditions of supporters (constructed from actions). Substitute increase and decrease with infinity assignments



- Relaxed Plan is  $\pi^+ = \{5 \times \text{water-plant}, \text{load-water}, 2 \times \text{east}, \text{south}\}$
- This plan solves Π<sup>+</sup>...does it solve Π, too?
- No, but its cost can be used to guide the search:

$$h^{MFF} = \sum_{a \in \pi^+} a^{\dagger}$$

- Computable?
  - Yes! it eventually terminates because of the mneed criteria
- Tractable?
  - Yes!  $\Pi^+$  can be decided in polytime
- Safe-pruning?
  - Yes! The existence of a plan for  $\Pi$  always implies that there is a relaxed plan for  $\Pi^+$
- Admissible?
  - No!

#### Intuition

Backward from the goal and recursively over subgoals that need to be achieved. Focus on subgoals independently



Two interpretations

- Pessimistic: sum cost of subgoals
- Optimistic: maximize cost of subgoals

#### Intuitively

An operation that combines some formula  $\psi$  with an action a to obtain a new formula  $\psi'$  which establishes sufficient and necessary conditions for reaching  $\psi$  through a

#### Formally

Regression is a function *regr* :  $\Gamma \times A \rightarrow \Gamma$  for which:

$$\forall \psi \in \mathsf{\Gamma}. \forall \mathsf{a} \in . \forall \mathsf{s} \in \mathsf{S}. \mathsf{s} \models \mathsf{regr}(\psi, \mathsf{a}) \iff \mathsf{s}[\mathsf{a}] \models \psi$$

- Regression is at the basis of subgoaling-based relaxation.
- In classical planning, if ψ is an atom and action a adds ψ, regr(ψ, a) = pre(a) - This is fundamental for making subgoaling so effective..

$$h^{max}(s,\psi) \doteq \begin{cases} 0 & \text{if } s \models \psi \\ \min_{a \in ach(s,\psi)} (h^{max}(s, \operatorname{pre}(a)) + c(a)) & \text{if } \psi \text{ is a PC} \\ \max_{\psi' \in \psi} (h^{max}(s,\psi')) & \text{if } \psi \text{ is } \wedge \end{cases}$$

- $ach(s, \psi)$  denotes the set of achievers of a proposition
- Decomposition comes from ignoring interactions among conjucts
- Pillar in optimal planning (Helmert and Domshlak 2006)

h<sup>max</sup> is:

- Admissible
- Safe-pruning
- Tractable

#### Question

How to extend it to consider numeric variables, but in particular numeric conditions?

#### Definition (Simple Numeric Condition)

A numeric condition is said to be simple if it is linear and all actions interacting with it are either increase or decrease by a constant

#### Definition (Simple Numeric Planning)

A numeric planning problem is said to be simple if it only contains simple numeric conditions.

## Regression in Simple Numeric Planning

#### Intuition

For each numeric formula, produce a new formula where each variable is substituted with the way such variable is modified by the action

#### Example

- $\psi: x + y \ge 10$ , an action *a* where eff(*a*) : x + = 1, and an action *b* where eff(*b*) : x = 1
- $regr(\psi, a) = (x+1) + y \ge 10$
- $\operatorname{regr}(\psi, b) = (x 1) + y \ge 10$

What can we say?

- Action a is a possible achiever for ψ. Also it seems you need at least 10 executions of a if you are in a state where both x and y are set to 0.
- Action b is not a possible achiever for  $\psi$ .
- These observations are independent on the state in which the actions are applied

How do we make this general?

### Possible achievers through multi-time regression

# Problem. Numeric effects are not idempotent, and we need a way to count the needed amount of actions

Observation:  $regr(...regr(\psi, a), a)$  m times is equivalent to operator  $\psi^{r(a,m)}$ :

$$\psi^{r(a,m)} \equiv m \underbrace{\left(\sum_{x=lhs(e), e \in eff_{num}(a)}^{N_{\psi,a}} w_{\psi,x} k_{a,x}\right)}_{x \in X} + \underbrace{\left(\sum_{x \in X} w_{\psi,x} x\right) + k_{\psi}}_{\xi \psi} \geq 0$$

This operator gives us:

- way to establish whether an action a is a possible achiever by inspecting the sign of  $N_{\psi,a}$ .
- easy counting mechanism:

$$\widehat{\operatorname{rep}}(a,\psi,s) = \begin{cases} 0 & \text{if } s \models \psi \\ \frac{-[\xi_{\psi}]^s}{N_{\psi,s}} & \text{else if } a \in \operatorname{ach}(\psi) \\ \infty & \text{otherwise} \end{cases}$$

$$h_{hbd}^{max}(s,\psi) \doteq \begin{cases} 0 & \text{if } s \models \psi \\ \min_{a \in ach(s,\psi)} (h_{hbd}^{max}(s, \operatorname{pre}(a)) + c(a)) & \text{if } \psi \in \operatorname{PC} \\ \min_{a \in ach(s,\psi)} (\widehat{\operatorname{rep}}(a,\psi,s) \cdot c(a) + h_{hbd}^{add}(\operatorname{pre}(a))) & \text{if } \psi \in \operatorname{SC} \\ \max_{\psi' \in \psi} h_{hbd}^{add}(s,\psi') & \text{if } \psi \text{ is } \wedge \end{cases}$$

Note that  $\widehat{rep}(\cdot, \cdot, \cdot)$  induces a continuous relaxation. The integral version of the minimisation problem is NP-Hard (Scala et al. 2020).

- Safe Pruning?
  - It can be shown by observing that one can always find all necessary achievers if unrelaxed problem is solvable
  - If heuristic yields infinity, the problem from that state is indeed unsolvable
- Tractable?
  - Can pose the problem as a blind search over atoms
  - Complexity polynomial on the number of actions and conditions of the problem
- Admissible?
  - NO!

#### Example Problem

- One goal x + y > 9, initial state is  $\{x = 0, y = 0\}$ .
- Two actions a and b such that
  - $pre(a) = (y \ge 10) eff(a) = \{x + = 1\}$
  - $pre(b) = \top eff(b) = \{(x+=0.5), (y+=5)\}$
- Cost optimal plan for this problem? 11! I.e.,  $\langle 2 \times b, 9 \times a 
  angle$
- Cost optimal through  $h_{hbd}^{max}$  is 12!
  - Obtained by choosing *a* as a best achiever, and recursively *b* for its precondition. Decomposition does not look into the positive interactions among the actions.
  - This is not a pathological case. Happens in real instances.

### Make it admissible

$$h_{hbd}^{max}(s,\psi) \doteq \begin{cases} 0 & \text{if } s \models \psi \\ \underset{a \in ach(s,\psi)}{\min} \left(h_{hbd}^{max}(s, \operatorname{pre}(a)) + c(a)\right) & \text{if } \psi \in \operatorname{PC} \\ \frac{\min}{a \in ach(s,\psi)} \left(\widehat{\operatorname{rep}}(a,\psi,s) \cdot c(a) + h_{hbd}^{add}(\operatorname{pre}(a))\right) & \text{if } \psi \in \operatorname{SC} \\ \underset{\psi' \in \psi}{\max} h_{hbd}^{add}(s,\psi') & \text{if } \psi \text{ is } \land \end{cases}$$

$$\left(\min_{a\in \mathsf{A}^+\cap ach(s,\psi)}\widehat{\operatorname{rep}}(a,\psi,s)\cdot c(a)\right) + \left(\min_{a\in \mathsf{A}^+\cap ach(s,\psi)}\hat{h}_{hbd}^{max}(s,\operatorname{pre}(a))\right)$$

### Search Tree with $h_{hbd}^{max}$ over a $10 \times 10$ Grid

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### Search Tree with no heuristic over a 10 $\times$ 10 Grid



$$h_{hbd}^{add}(s,\psi) \doteq \begin{cases} 0 & \text{if } s \models \psi \\ \min_{a \in ach(s,\psi)} \left( h_{hbd}^{add}(s, \operatorname{pre}(a)) + c(a) \right) & \text{if } \psi \in \mathsf{PC} \\ \min_{a \in ach(s,\psi)} \left( \widehat{\operatorname{rep}}(a,\psi,s) \cdot \gamma(a) + h_{hbd}^{add}(\operatorname{pre}(a)) \right) & \text{if } \psi \in \mathsf{SC} \\ \sum_{\psi' \in \psi} h_{hbd}^{add}(s,\psi') & \text{if } \psi \text{ is } \land \end{cases}$$

- Pretend all subproblems don't interact. That is, sum the cost of each subproblem
- Can work much (MUCH) better than admissible variant in practice.
- Lots of research tries to understand when things can be summed up or not. Underlying problem is NP-Hard.

- It depends on the problem!
- Subgoaling enables optimal planning. There are at the moment no heuristics based on intervals that do so (Piacentini et al. 2018 is a slightly different exception). Interval-based relaxation with cost limited, but something exists (Aldinger et al. 2015)
- Subgoaling is difficult to be extended with more general expressions, action effects.

- Want to build a numeric planner?
  - Can use state-space search onto which you can use best-first search schema pretty much off the shelf.
- Heuristic reasoning is crucial. For this, two different schemata are known: based on Subgoals decomposition or Variables decomposition.
- Both have strengths and weaknesses. Be careful to understand your problem better first. Simplify if complexity is not necessary!
- Plenty of problems in combining the two relaxations (Piacentini et al 2018), and extend to more powerful constructs (e.g., non-linear problems, processes...)

- LPRPG heuristic. Builds on intervals, but reason over resource-flows patterns much more effectively than what we have seen here (Coles et al. 2008)
- LPG planner uses different variants of relaxed plan heuristics that are aimed at capturing some action interference (Gerevini et al. 2008). The SAPA planner does something similar (Do and Kambhabati 2003)
- Using Linear programs to devise optimal estimates (Piacentini et al. 2018)
- Better integration of the heuristic in the search through helpful actions (Hoffmann 2003)
- Linear programs also in variants of  $h^{max}$  with the aim of capturing conflicts among subgoals (Scala et al. 2020)

### Non-exhaustive List of Available Numeric Planners

- Optic, Colin (Kings College). PDDL 2.1, Fragments of PDDL+. No discretisation, Heuristic search via IBR. C++
- Metric-FF (Joerg Hoffmann). PDDL 2.1. Heuristic from IBR. Restricted to linear and acyclic tasks, C
- LPG (University of Brescia). PDDL 2.1, monotone in each variable separately. Local Search with IBR heuristic. C
- ENHSP (Enrico Scala, ANU, University of Brescia). PDDL2.1, PDDL+, optimality guarantees, highly customizable. Heuristic search via IBR, Subgoaling. JAVA
- SMTPLAN (King's College). Fragment of PDDL+. Depends on SMT engine. Frotend in Python
- TFD (G. Roger, P. Eyerich, C. Dornhege, R. Matmuller) PDDL2.1, heuristic search through variant of *h*<sup>hadd</sup> heuristic, C++

#### All can be found easily by googling them

# **Main References for Numeric Planning**

Modeling Language and Complexity Analysis

Fox, Maria, and Derek Long. "PDDL2.1: An extension to PDDL for expressing temporal planning domains." *Journal of artificial intelligence research* 20 (2003): 61-124.

Helmert, Malte. "Decidability and Undecidability Results for Planning with Numerical State Variables." In *AIPS*, pp. 44-53. 2002.

Haslum, Patrik, Nir Lipovetzky, Daniele Magazzeni, and Christian Muise. "An introduction to the planning domain definition language." *Synthesis Lectures on Artificial Intelligence and Machine Learning* 13, no. 2 (2019): 1-187.

Geffner, Hector, and Blai Bonet. "A concise introduction to models and methods for automated planning." Synthesis Lectures on Artificial Intelligence and Machine Learning 8, no. 1 (2013): 1-141.

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Piacentini, Chiara, Maria Fox, and Derek Long. "Planning with numeric timed initial fluents." In *Proceedings of the Twenty-Ninth AAAI Conference on Artificial Intelligence*, pp. 4196-4197. 2015. Coles, Amanda, M. Fox, and D. Long. "A hybrid LP-RPG heuristic for modelling numeric resource flows in planning." *Journal of Artificial Intelligence Research* 46 (2013): 343-412.

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Gerevini, Alfonso E., Alessandro Saetti, and Ivan Serina. "An approach to efficient planning with numerical fluents and multi-criteria plan quality." *Artificial Intelligence* 172, no. 8-9 (2008): 899-944. Do, Minh, and Subbarao Kambhampati. "Sapa: A multi-objective metric temporal planner." Journal of Artificial Intelligence Research 20 (2003): 155-194.

Mainly Used Available Planners (alphabetical order)

ENHSP - <u>https://sites.google.com/view/enhsp/</u> E. Scala (see website for other contributors to the platform)

LPG - https://lpg.unibs.it/lpg/ A.E. Gerevini, A. Saetti, I. Serina

LPRPG - https://nms.kcl.ac.uk/planning/software/lprpg.html Amanda Coles, Andrew Coles, M. Fox, D. Long

Metric-FF - https://fai.cs.uni-saarland.de/hoffmann/metric-ff.html J. Hoffmann

Optic - <u>https://nms.kcl.ac.uk/planning/software/optic.html</u> Amanda Coles, Andrew Coles, M. Fox, D. Long

SMTPLAN - https://github.com/KCL-Planning/SMTPlan M. Cashmore, D. Magazzeni, M. Fox, D. Long TFD - <u>http://gki.informatik.uni-freiburg.de/tools/tfd/</u> G. Röger, P. Eyerich, C. Dornhege, and R. Mattmüller.

UPMURPHI - https://github.com/gdellapenna/UPMurphi G. Della Penna, D. Magazzeni, F. Mercorio